

Introduction to Natural Language Processing

Part VI: NLP using Regular Expressions

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Learning Objectives

Concepts

- Different types of formal grammars in NLP
- Sequential patterns in language
- Main operators of regular expressions

Methods

- Identification of numeric entities with regular expressions

Covered tasks

- Time expression recognition

Outline of the Course

- I. Overview
- II. Basics of Linguistics
- III. NLP using Rules
- IV. NLP using Lexicons
- V. Basics of Empirical Methods
- VI. NLP using Regular Expressions
 - Introduction
 - Regular Grammars
 - Regular Expressions in NLP
- VII. NLP using Context-Free Grammars
- VIII. NLP using Language Models
- IX. Practical Issues

Introduction

Woodchuuucks

HOW MUCH WOOD WOULD A **WOODCHUCK** CHUCK,
IF A **WOODCHUCK** COULD CHUCK WOOD?

So much wood as a **woodchuck** could,
if a **woodchuck** would chuck wood.

It would chuck, if **it** would, as much wood as
it could, if a **woodchuck** could chuck wood.

A **woodchuck** would chuck no amount of wood,
since **woodchucks** can't chuck wood.

*But if **Woodchucks** Could and Would Chuck Some Wood,
What Amount of Wood Would each **Woodchuck** Chuck?*

... and so forth

If you don't know, maybe you know the difference between
woodchucks and **groundhogs**?

There is none. A **groundhog** is a **woodchuck**.



<https://commons.wikimedia.org>

Woodchuuucks

Mining Woodchucks from Text

Woodchucks in the text above

- “Woodchuuucks”
- “WOODCHUCK”
- “woodchuck”
- “woodchucks”
- “Woodchucks”
- “Woodchuck”
- “groundhogs”
- “groundhog”

along with “It” and “it”



Questions

- How to find these and other references to woodchucks automatically?
- What makes woodchuck mining challenging?

Examples: Ambiguous cases, missing whitespaces, varying writings, ...

Contact Information

Finding contact information

- See the following e-mail:



Henning Wachsmuth

Regular expressions

An: Henning Wachsmuth

Dear students,

how does the mail tool infer that the following lines contain contact information:

Marty McFly
9303 Lyon Drive
Hill Valley, CA 95420
USA
marty.mcfly@delorian.bttf ✓

Does this have to do anything with regular expressions?



<https://pixabay.com>

Questions

- How can we describe patterns in text sequences formally?
- What can be described well, what not?

NLP using Regular Expressions

Regular expression (regex)

- A sequence of characters that describes sequential text patterns
- A text can be matched against a regex to find all pattern instances

Use in NLP

- Effective for finding information that follows clear sequential structures
- **Examples.** Numeric entities, structural entities (e.g., e-mail addresses), lexico-syntactic relations (e.g., “<NN> is a <NN>”), ...

Numeric (and alphanumeric) entities

- Values, quantities, proportions, ranges, or similar
- This includes time periods, dates, monetary values, phone numbers, ...

“in this year”

“2024-06-13”

“\$ 100 000”

“762-123 77”

Numeric entity recognition

- The text analysis that mines numeric entities from text
- Used in NLP within many information extraction tasks

Regular Grammars

Grammar

Grammar.

The difference between
knowing your shit and
knowing you're shit.



<https://someecards.com>

Grammar

Grammars

Grammars

- A grammar is a description of the valid structures of a language.
- One of the most central concepts of linguistics is *formal grammars*.

Formal grammars

- A formal grammar specifies a set of rules consisting of terminal and non-terminal symbols.
- **Terminals.** “Words” that cannot be rewritten any further
- **Non-terminals.** Clusters or generalizations of terminals

Grammar (Σ, N, S, R)

Σ An alphabet, i.e., a finite set of terminal symbols, $\Sigma = \{v_1, v_2, \dots\}$

N A finite set of non-terminal symbols, $N = \{W_1, W_2, \dots\}$

S A start non-terminal symbol, $S \in N$

R A finite set of production rules, $R \subseteq (\Sigma \cup N)^+ \setminus \Sigma^* \times (\Sigma \cup N)^*$

Here, $\Sigma^* := \Sigma \times \Sigma \times \dots$

Grammar

Chomsky Grammars

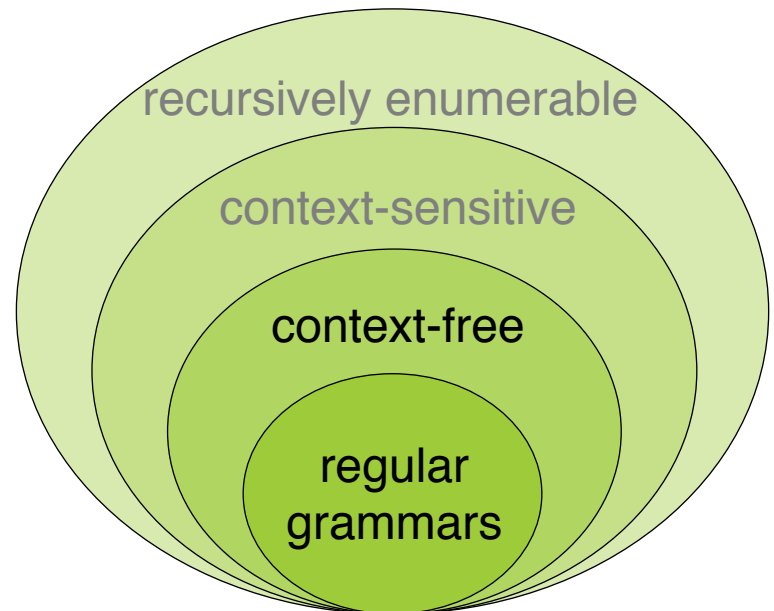
Four types of formal grammars

- Chomsky-0 (recursively enumerable). Any (Σ, N, S, R) as defined
- Chomsky-1 (context-sensitive). Only rules $U \rightarrow V$ with $|U| \leq |V|$
- Chomsky-2 (context-free). Only rules $U \rightarrow V$ with $U \in N$
- Chomsky-3 (regular). Only rules $U \rightarrow V$ with $U \in N$ and $V \in \{\varepsilon, v, vW\}$,
 $v \in \Sigma, W \in N$

ε is the empty word.

Grammars in NLP

- Only regular and context-free grammars are commonly used.
- Regular grammars are covered here, context-free grammars in the next part.



Regular Grammars

Regular grammar

- A grammar (Σ, N, S, R) where all rules in R are of the form $U \rightarrow V$ with $U \in N$ and $V \in \{\varepsilon, v, vW\}$, such that $v \in \Sigma$ and $W \in N$
- **Extended.** A regular grammar is *extended*, if $v \in \Sigma^*$
Below, we refer to this also as *regular grammar* only.
- **Right-regular.** Intuitively, a structure defined by a regular grammar can be constructed from left to right.

Regular language

- A language is regular, if there is a regular grammar that defines it.

Representation

- Every regular grammar can be represented by a *finite-state automaton*.
- Every regular grammar can be represented by a *regular expression*.

And vice versa

Regular Grammars

Finite-State Automata

Finite-state automaton (FSA) (recap)

- An FSA is a state machine that reads a string from a regular language.
- It represents the set of all strings belonging to the language.

FSA as a 5-tuple $(Q, \Sigma, q_0, F, \delta)$

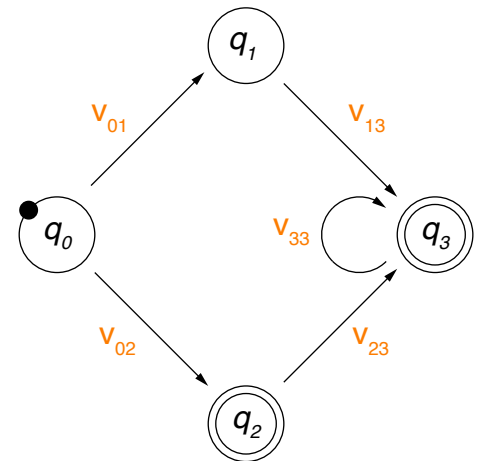
Q A finite set of $n > 0$ states, $Q = \{q_0, \dots, q_n\}$

Σ An alphabet, i.e., a finite set of terminal symbols, $\Sigma = \{v_1, v_2, \dots\}$, $\Sigma \cap Q = \emptyset$

q_0 A start state, $q_0 \in Q$

F A set of final states, $F \subseteq Q$

δ A transition function between states, triggered based on $v \in \Sigma$, $\delta : Q \times \Sigma \rightarrow Q$



Regular Expressions in NLP

Regular Expressions

Regular expression (regex)

- A regex defines a regular language over an alphabet Σ as a sequence of characters from Σ and *metacharacters*
- Metacharacters denote disjunction, negation, repetition, ... (see below)
- **Example.** The FSA above is defined by the following regex:

$$v_{02} \mid (v_{01}v_{13} \mid v_{02}v_{23}) v_{33}^*$$

Use of regular expressions

- Definition of patterns that generalize over structures of a language
- A pattern matches all spans of text that contain any of the structures.

Regular expressions in NLP

- Regexes are a widely used technique in NLP, particularly for the extraction of numeric and similar entities.
- In statistical NLP, regexes often take on the role of features.

Regular Expressions

Characters and Metacharacters

Regular characters

- The default interpretation of a character sequence in a regex is a concatenation of each single character.

woodchuck matches “woodchuck”

Metacharacters

- A regex uses specific characters to encode specific regular language constructions, such as negation and repetition.
- The main metacharacters are the following (in Python notation):

[] - | ^ . () \ * + ?

The characters used partly differ across literature and programming languages.

- Some languages also include certain *non-regular* constructions, e.g.,
`\b` matches if a word boundary is reached.

Regexes can solve this case only if given token information.

Regular Expressions

Disjunction

Disjunction of patterns

- Brackets `[]` specify a character class.

`[wW]` matches “w” or “W”

`[wod]` matches “w” or “o” or “d”

- Disjunctive ranges of characters can be specified with a hyphen `-`.

`[a-zA-Z]` matches any letter

`[0-8]` matches any digit except for “9”

- The pipe `|` specifies a disjunction of string sequences.

`groundhog|woodchuck` matches “groundhog” and “woodchuck”

Notes on disjunctions

- Different disjunctions can be combined.

`[gG]roundhog|[wW]oodchuck` matches “groundhog”, “Woodchuck”, ...

- In Python, many metacharacters are not active within brackets.

`[wod.]` matches “w”, “o”, “d”, and “.”

Regular Expressions

Negation, Choice, Grouping

Negation

- The caret `^` inside brackets complements the specified character class.

`[^0-9]` matches anything but digits

`[^wo]` matches any character but “w”, “o”

Free choice

- The period `.` matches any character.

To match a period, it needs to be escaped as: `\.`

`w..dchuck` matches “woodchuck”, “woudchuck”, ...

Grouping

- Parentheses `()` can be used to group parts of a regex. A grouped part is treated as a single character.

`w(oo|uo|ou)dchuck` matches “woodchuck”, “wodchuck”, and “wodchuck”

Regular Expressions

Whitespaces and Predefined Character Classes

Whitespaces

- Different whitespaces are referred to with different special characters.
- For instance, `\n` is the regular new-line space.

Predefined character classes

- Several specific character classes are referred to by a backslash `\` followed by a specific letter.

<code>\d</code>	Any decimal digit	equivalent to <code>[0-9]</code>
<code>\D</code>	Any non-digit character	equivalent to <code>[^0-9]</code>
<code>\s</code>	Any whitespace character	equivalent to <code>[\t\n\r\f\v]</code>
<code>\S</code>	Any non-whitespace character	equivalent to <code>[^\t\n\r\f\v]</code>
<code>\w</code>	Any alphanumeric character	equivalent to <code>[a-zA-Z0-9_]</code>
<code>\W</code>	Any non-alphanumeric character	equivalent to <code>[^a-zA-Z0-9_]</code>

- These classes can be used within brackets.

`[\s0-9]` matches any space and digit.

Regular Expressions

Repetition

Repetition

- The asterisk `*` repeats the previous character zero or more times.

`woo*dchuck` matches “`wodchuck`”, “`woodchuck`”, “`woodchuck`”, ...

- The plus `+` repeats the previous character one or more times.

`woodchu+ck` matches “`woodchuck`”, “`woodchuuck`”, “`woodchuuuck`”, ...

- The question mark `?` repeats the previous character zero or one time.

`woodchucks?` matches “`woodchuck`” and “`woodchucks`”

Notes on repetitions

- Repetitions are implemented in a greedy manner in many programming languages, i.e., longer matches are preferred over shorter ones.

`to*` matches “`too`”, not “`too`”, ...

- This may actually violate the regularity of the defined language.

“`woodchuck`” needs to be processed twice for the regex `wo*odchuck`

Regular Expressions

Summary of Metacharacters

Char	Concept	Example
[]	Disjunction of characters	<code>[Ww]oodchuck</code>
–	Ranges in disjunctions	There are <code>[0–9]+</code> woodchucks\.
	Disjunction of regexes	<code>woodchuck groundhog</code>
^	Negation	<code>[^0–9]</code>
.	Free choice	What a <code>(.)*</code> woodchuck
()	Grouping of regex parts	w <code>(oo)+</code> dchuck
\	Special (sets of) characters	<code>\s</code> woodchuck <code>\s</code>
*	Zero or more repetitions	woo <code>o*</code> dchuck
+	One or more repetitions	woodchu <code>u+</code> ck
?	Zero or one repetition	woodchuck <code>s?</code>

Regular Expressions

Examples

The

- Regex for all variations of “the” in news article text:

`the` (misses capitalized cases, matches “theology”, ...)

`[^a-zA-Z][tT]he[^a-zA-Z]` (requires a character before and afterwards)

Woodchucks

- Regex for all woodchuck cases from above (and for similar):

`( [wW][oO][oO][dD][cC][hH][uU]+[cC][kK] | [Gg]roundhog ) [sS]?`

E-mail addresses

- All e-mail addresses with the German top-level domain, whose text segments contain no special characters:

`[a-zA-Z0-9]+ @ ([a-zA-Z0-9]+\.) * [a-zA-Z0-9][a-zA-Z0-9]+ \.de`

Time Expression Recognition with Regular Expressions

Time expression

- An alphanumeric entity that represents a date or a period

“Cairo, **August 25th 2010** — Forecast on Egyptian Automobile industry
[...] **In the next five years**, revenues will rise by 97% to US-\$ 19.6 bn. [...]”

Time expression recognition

- The text analysis that finds time expressions in natural language text
- Used in NLP within temporal relation extraction and event extraction

Approach in a nutshell

- Model sequential structure of time expressions with a complex regex.
- Include lexicons derived from training data to match closed-class terms.
Example closed classes: Month names, prepositions, ...
- Match regex with sentences of a text.

Notice

- The approach can easily be adapted to other types of information.

Time Expression Recognition with Regular Expressions

Pseudocode

Signature

- **Input.** A text split into sentences, and a regex
- **Output.** All time expressions in the text

extractAllMatches (**List**<**Sentence**> sentences, **Regex** regex)

```
1.      List<TimeExpression> matches ← ()
2.      for each sentence ∈ sentences do
3.          int index ← 0
4.          while index < sentence.length - 1 do
5.              int [] exp ← regex.match(sentence.sub(index))
6.              if exp ≠ ⊥ then    // ⊥ represents "null"
7.                  matches.add(new TimeExpression(exp[0], exp[1]))
8.                  index ← exp[1]
9.              index ← index + 1
10.     return matches
```

Notice

- Most programming languages provide explicit matching classes.

Time Expression Recognition with Regular Expressions

Complete Regex (Part 1 out of 2)

```

((([iI]n|[wW]ithin|[tT]o|s|s?he| [tT]o| [fF]or|s|s?the| [fF]or| [fF]rom|[sS]ince|[aA]fter|[bB]efore|[bB]etween|[aA]t|[oO]n|[oO]ver|[pP]er)((\s+(\r(\n)?|\n)?|(\r(\n)?|\n)))*s*(([tT]he|[tT]his|[tT]hese|[tT]hose|[iI]ts)?)(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*))?((([123456789])|[12]|d|3[01])(\.(\/))(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*))?((month|time|span)?(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*(from(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*))?(([J]anuary|[J]an\.[J]an|[F]ebruary|[F]eb\.[F]eb|[M]arch|[M]ar\.[M]ar|[A]pril|[A]pr\.[A]pr|[M]ay|[J]une|[J]un\.[J]un|[J]uly|[J]ul\.[J]ul|[A]ugust|[A]ug\.[A]ug|[S]eptember|[S]ep\.[S]ep|[O]ctober|[O]ct\.[O]ct|[N]ovember|[N]ov\.[N]ov|[D]ecember|[D]ec\.[D]ec|[S]pring|[S]ummer|[A]utumn|[F]all|[W]inter))((([123456789])|[1012])(\.(\/))((([19]20)?d2)?|(((([iI]n|[wW]ithin|[tT]o|s|s?the|[tT]o|[fF]or|s|s?the|[fF]or|[fF]rom|[sS]ince|[aA]fter|[bB]efore|[bB]etween|[aA]t|[oO]n|[oO]ver|[pP]er)((\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*))s*(([tT]he|[tT]his|[tT]hese|[tT]hose|[iI]ts)?)(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*((([S]tart|[B]egin|[S]tart|[B]egin|[E]nd|[E]nd|[M]idth|[M]idth)(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*((([S]tart|[B]egin|[S]tart|[B]egin|[E]nd|[E]nd|[M]idth|[M]idth)(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*(([tT]he|[tT]his|[tT]hese|[tT]hose|[iI]ts)?)(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*(([a-z]]+(\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*))?((([lL]ast|[pP]receding|[pP]ast|[C]urrent|[tT]his|[uU]pcoming|[fF]ollowing|[sS]ucceeding|[nN]ext))((\s+(\r(\n)?|\n)?|(\r(\n)?|\n))\s*((([12]3[4]5[6]7[8]9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|3|4|5|6|7|8|9)d?|([oO]ne|[sS]everal|[sS]ome|[bB]oth|[tT]wo|[tT]hree|[fF]our|[fF]ive|[sS]ix|[sS]even|[eE]ight|[nN]ine|[tT]en|[eE]leven|[tT]welve|[tT]wenty|[tT]hirty|[fF]ourty|[fF]ifty|[sS]ixty|[sS]eventy|[eE]ighty|[nN]inety|[hH]undred|[aA]\s?s?hundred))((([1012]?|2|3|4|5|6|7|8|9)(\.(\/))|([fF]irst|[sS]econd|[tT]hird|[fF]ourth|[fF]ifth|[sS]ixth|[sS]eventh|[eE]ighth|[nN]inth|[tT]enth|[eE]leventh))(-(([1012]?|2|
```

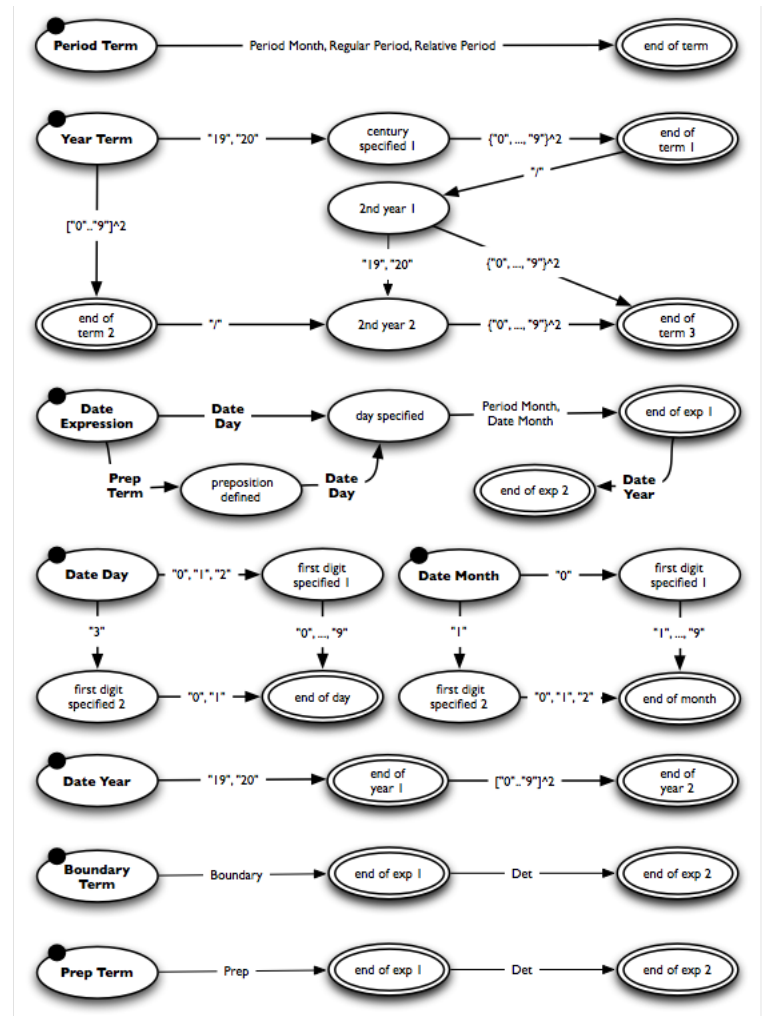
Time Expression Recognition with Regular Expressions

Complete Regex (Part 2 out of 2)

[A] [S] [S] [hundred]))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s*) | ((((1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) \d ? | ([oO] ne | [sS] everal | [sS] ome | [bB] oth | [tT] wo | [tT] hree | [fF] our | [fF] ive | [sS] ix | [sS] even | [eE] ight | [nN] ine | [tT] en | [eE] leven | [tT] welve | [tT] wenty | [tT] thirty | [fF] ourty | [fF] orte | [fF] ifty | [fF] ifty | [sS] ixty | [sS] eventy | [eE] ighty | [nN] inety | [hH] undred | [aA] \s?s?hundred)) | ((1 | 0 | 12 ? | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) (\. | ()) | ([fF] irst | [sS] econd | [tT] hird | [fF] ourth | [fF] ifth | [sS] ixth | [sS] eventh | [eE] ighth | [nN] inth | [tT] enth | [eE] leventh))) - ((1 | 0 | 12 ? | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) (\. | ()) | ([fF] irst | [sS] econd | [tT] hird | [fF] ourth | [fF] ifth | [sS] ixth | [sS] eventh | [eE] ighth | [nN] inth | [tT] enth | [eE] leventh))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ((1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) \d ? | ([oO] ne | [sS] everal | [sS] ome | [bB] oth | [tT] wo | [tT] hree | [fF] our | [fF] ive | [sS] ix | [sS] even | [eE] ight | [nN] ine | [tT] en | [eE] leven | [tT] welve | [tT] thirty | [fF] ourty | [fF] ifty | [fF] orte | [fF] ifty | [sS] ixty | [sS] eventy | [eE] ighty | [nN] inety | [hH] undred | [aA] \s?s?hundred))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* (([lL] ast | [pP] receding | [pP] last | [cC] urrent | [tT] his | [uU] pcoming | [fF] ollowing | [sS] uceeding | [nN] ext))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s*) ? ((Q (1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) \d ? | (1 | 2) (\w ([a-z]) * \s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s*) ? (year | quarter) ([a-z]) *) | (month | time (span) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* (from (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s*) ? ([Jj] anuary | [Jj] an\ . | [Jj] an | [Ff] ebruary | [Ff] eb\ . | [Ff] eb | [Mm] arch | [Mm] ar\ . | [Mm] ar | [Aa] pril | [Aa] pr\ . | [Aa] pr | [Mm] ay | [Jj] une | [Jj] un\ . | [Jj] un | [Jj] uly | [Jj] ul\ . | [Jj] ul | [Aa] ugust | [Aa] u g\ . | [Aa] ug | [Ss] eptember | [Ss] ep\ . | [Ss] ep | [Oo] ctober | [Oo] ct\ . | [Oo] ct | [Nn] ovember | [Nn] ov\ . | [Nn] ov | [Dd] ecember | [Dd] ez\ . | [Dd] ez | [Ss] pring | [Ss] ummer | [Aa] utumn | [Ff] all | [Ww] inter)) | (([Rr] eported\s?s?time\s?s?span | [Rr] eported\s?s?time\s?s?span | [Rr] eported\s?s?time\s?s?span | [Rr] eported\s?s?time | [rR] eported\s?s?time | [Tt] ime\s?s?span | [Tt] ime\s?s?span | [Tt] ime\s?s?span | [Ss] pan | [Ss] pan | [Dd] ecade | [Dd] ecade))) (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ((19 | 20) \d2 / (19 | 20) ? \d2) ? | (19 | 20) \d2 / (19 | 20) ? \d2) ? | (19 | 20) \d2 / (19 | 20) ? \d2) ? | (19 | 20) \d2 / (19 | 20) ? \d2) ?)) (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ([tT] o | [aA] nd | [oO] r | [oO] n | [aA] t | [oO] f\s?s?the | [oO] f | [tT] he | [tT] his | [iI] t | [iI] nstead\s?s?of ((\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* (([Ss] tart | [bB] egin | [Ss] tart | [Bb] egin | [Ee] nd | [eE] nd | [Mm] idth | [mM] idth) (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ([tT] he | [tT] his | [tT] hese | [tT] hose | [iI] ts))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* [a-z] +) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ((([lL] ast | [pP] receding | [pP] ast | [cC] urrent | [tT] his | [uU] pcoming | [fF] ollowing | [sS] uceeding | [nN] ext))) (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ((1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) \d ? | ([oO] ne | [sS] everal | [sS] ome | [bB] oth | [tT] wo | [tT] hree | [fF] our | [fF] ive | [sS] ix | [sS] even | [eE] ight | [nN] ine | [tT] en | [eE] leven | [tT] welve | [tT] thirty | [fF] ourty | [fF] ifty | [fF] orte | [fF] ifty | [sS] ixty | [sS] eventy | [eE] ighty | [nN] inety | [hH] undred | [aA] \s?s?hundred)) | ((1 | 0 | 12 ? | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) (\. | ()) | ([fF] irst | [sS] econd | [tT] hird | [fF] ourth | [fF] ifth | [sS] ixth | [sS] eventh | [eE] ighth | [nN] inth | [tT] enth | [eE] leventh)) - ((1 | 0 | 12 ? | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) (\. | ()) | ([fF] irst | [sS] econd | [tT] hird | [fF] ourth | [fF] ifth | [sS] ixth | [sS] eventh | [eE] ighth | [nN] inth | [tT] enth | [eE] leventh))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ((1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) \d ? | ([oO] ne | [sS] everal | [sS] ome | [bB] oth | [tT] wo | [tT] hree | [fF] our | [fF] ive | [sS] ix | [sS] even | [eE] ight | [nN] ine | [tT] en | [eE] leven | [tT] welve | [tT] thirty | [fF] ourty | [fF] ifty | [fF] orte | [fF] ifty | [sS] ixty | [sS] eventy | [eE] ighty | [nN] inety | [hH] undred | [aA] \s?s?hundred))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* (([lL] ast | [pP] receding | [pP] ast | [cC] urrent | [tT] his | [uU] pcoming | [fF] ollowing | [sS] uceeding | [nN] ext))) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s*) ? ((Q (1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9) \d2 ? | (1 | 2) (\w ([a-z]) * \s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s*) ? (year | quarter) ([a-z]) *) | (month | time (span) ? (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* (from (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s*) ? ([Jj] anuary | [Jj] an\ . | [Jj] an | [Ff] ebruary | [Ff] eb\ . | [Ff] eb | [Mm] arch | [Mm] ar\ . | [Mm] ar | [Aa] pril | [Aa] pr\ . | [Aa] pr | [Mm] ay | [Jj] une | [Jj] un\ . | [Jj] un | [Jj] uly | [Jj] ul\ . | [Jj] ul | [Aa] ugust | [Aa] ug\ . | [Aa] ug | [Ss] eptember | [Ss] ep\ . | [Ss] ep | [Oo] ctober | [Oo] ct\ . | [Oo] ct | [Nn] ovember | [Nn] ov\ . | [Nn] ov | [Dd] ecember | [Dd] ez\ . | [Dd] ez | [Ss] pring | [Ss] ummer | [Aa] utumn | [Ff] all | [Ww] inter)) | (([Rr] eported\s?s?time\s?s?span | [Rr] eported\s?s?time\s?s?span | [Rr] eported\s?s?time | [rR] eported\s?s?time | [Tt] ime\s?s?span | [Tt] ime\s?s?span | [Tt] ime\s?s?span | [Ss] pan | [Ss] pan | [Dd] ecade | [Dd] ecade))) (\s+ (\r (\n) ? | \n) ? | (\r (\n) ? | \n)) \s* ((19 | 20) \d2 / (19 | 20) ? \d2) ? | (19 | 20) \d2 / (19 | 20) ? \d2) ? | (19 | 20) \d2 / (19 | 20) ? \d2) ? | (19 | 20) \d2 / (19 | 20) ? \d2) ?)) *

Time Expression Recognition with Regular Expressions

Complete Regex as a Finite-State Automaton



Time Expression Recognition with Regular Expressions

Top-level FSA of Complete Regex



Notice

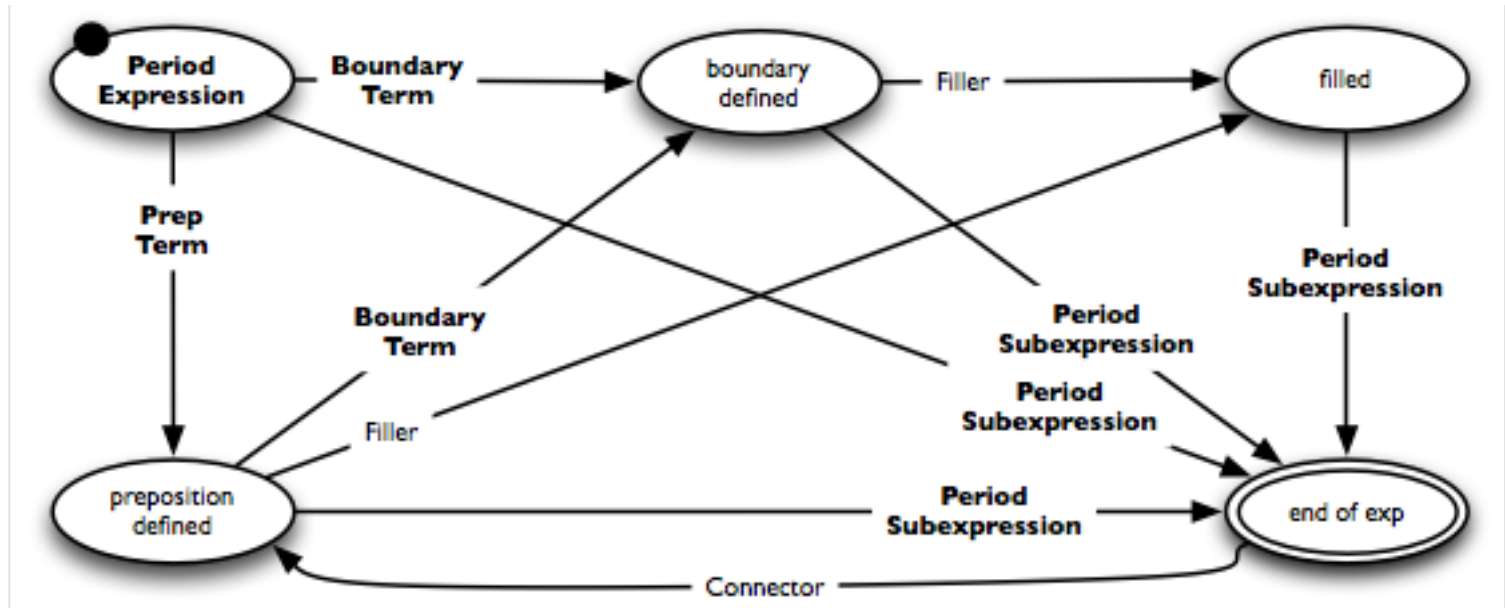
- Bold edge labels indicate sub-FSAs, regular labels indicate lexicons.
- Below, the FSA of period expressions is decomposed top-down.
The regex for date expressions is left out for brevity.
- During development, building a regex rather works bottom-up.

Example period expression

- “From the very end of last year to the 2nd half of 2019”
prep filler boundary relational period connector ordinal period year

Time Expression Recognition with Regular Expressions

Sub-FSA for Period Expressions

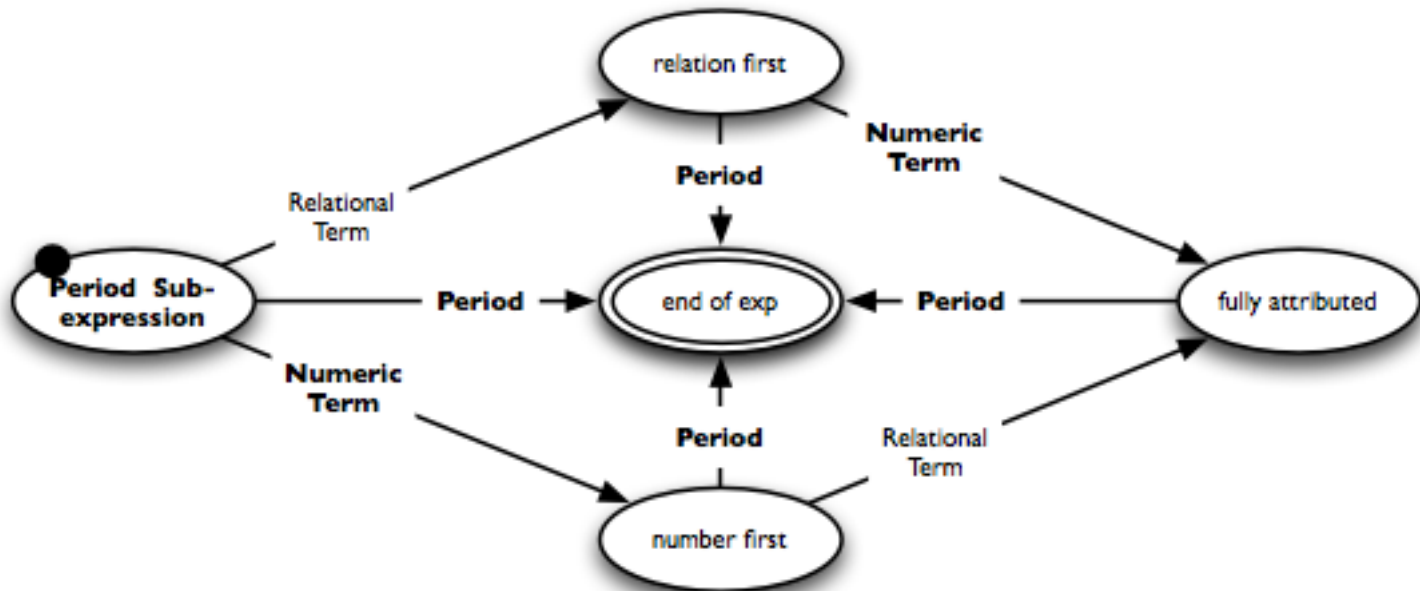


Lexicons

- **Connector lexicon.** “to the”, “to”, “and”, “of the”, “of”, ...
- **Fillers.** Any single word, such as “**very**” in the example above

Time Expression Recognition with Regular Expressions

Sub-FSA for Period Subexpressions

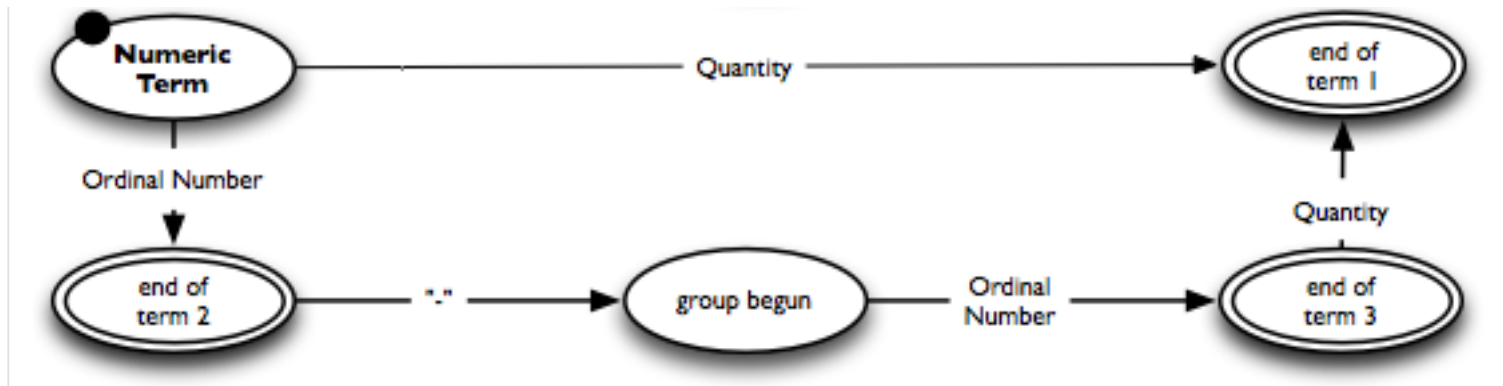


Lexicons

- **Relational term lexicon.** “last”, “preceding”, “past”, “current”, “this”, “upcoming”, “next”, ...

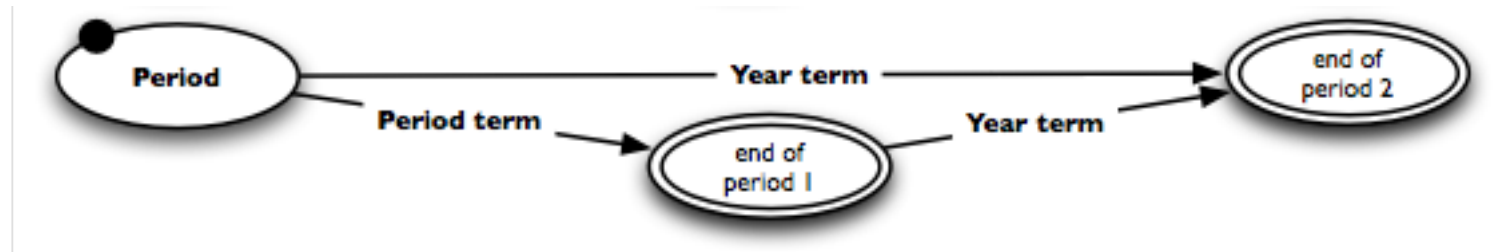
Time Expression Recognition with Regular Expressions

Sub-FSAs for Numeric Terms and Periods



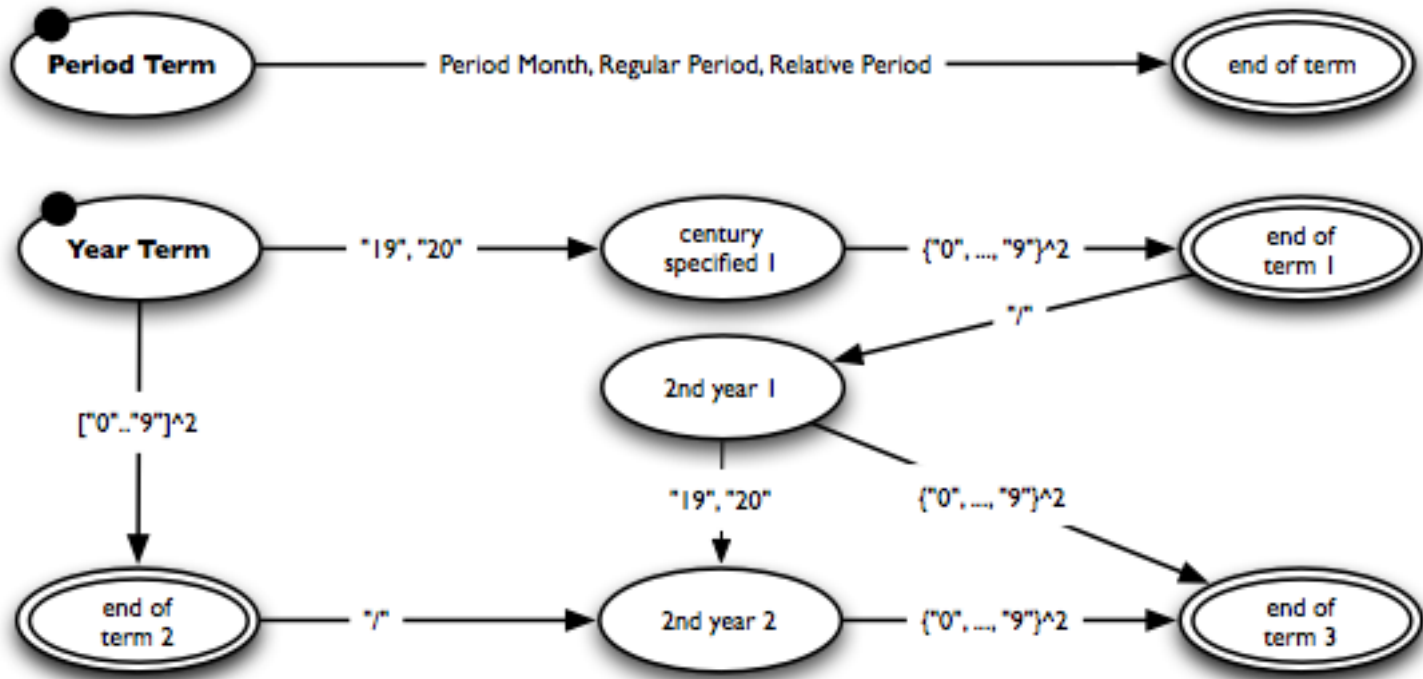
Lexicons

- Ordinal number lexicon. “first”, “1st”, “second”, “2nd”, “third”, “3rd”, ...
- Quantity lexicon. “one”, “two”, “three”, “both”, “several”, “a hundred”, ...



Time Expression Recognition with Regular Expressions

Sub-FSAs for Period Terms and Year Terms

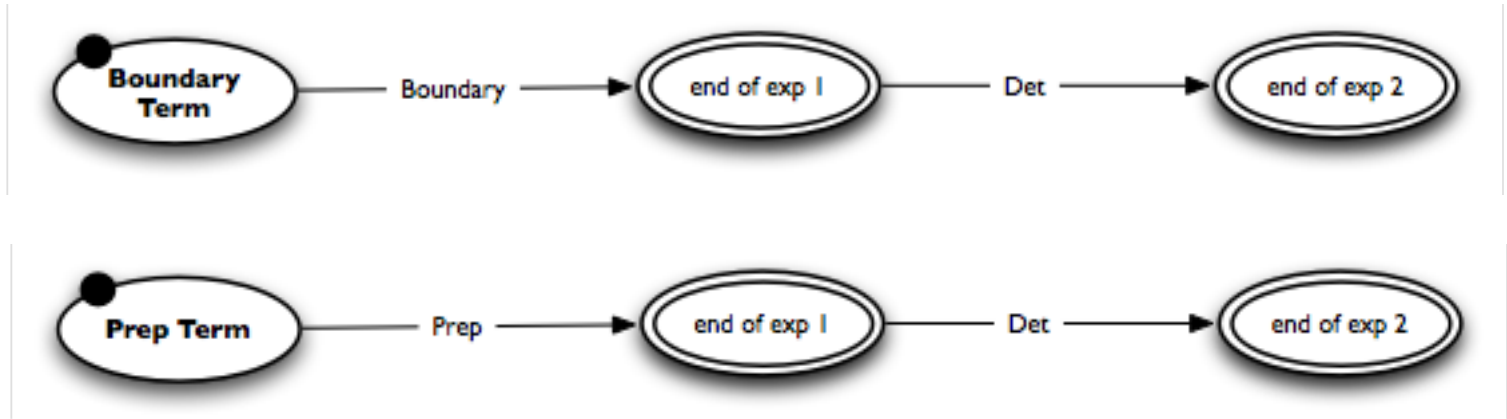


Lexicon

- Period month lexicon. "March", "Mar.", "Mar", "Fall", "fall", "Autumn", ...
- Regular period lexicon. "year", "month", "quarter", "half", ...
- Relative period lexicon. "decade", "reported time", "time span", ...

Time Expression Recognition with Regular Expressions

Sub-FSAs for Boundary Terms and Prepositional Terms



Lexicons

- **Boundary lexicon.** “Beginning”, “beginning”, “End”, “end”, “Midth”, ...
- **Prep lexicon.** “in”, “within”, “to”, “for”, “from”, “since”, ...
- **Det lexicon.** “the”, “a”, “an”

Time Expression Recognition with Regular Expressions

Evaluation

Effectiveness of the time expressions (Wachsmuth, 2015)

- Originally developed for German texts; only this version was evaluated
- **Data.** Test set of *InfexBA Corpus* with 6038 business news sentences
- **Evaluation metrics.** Precision, recall, F_1 -score, run-time per sentence

Run-time measured on a standard computer from 2009

Results compared to other analyses

Task	Technique	Precision	Recall	F_1 -score	ms/sent.
Time expression recognition	Regex	0.91	0.97	0.94	0.36
Money expression recognition	Regex	0.99	0.95	0.97	0.68
Forecast event extraction	Linear SVM	0.87	0.93	0.90	0.81
Time/Money relation extraction	Linear SVM	0.75	0.88	0.81	10.41

Observations

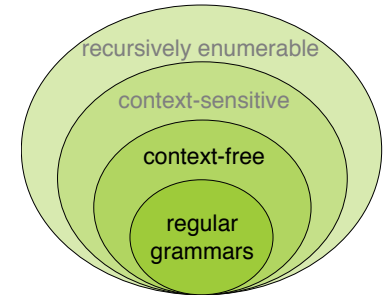
- Regexes for semi-closed-class entity types such as time expressions can achieve very high effectiveness and efficiency.
- Their development is complex and time-intensive, though.

Conclusion

Conclusion

Formal grammars

- Descriptions of valid structures in a language
- Production rules map non-terminals to terminals
- Regular grammars model sequential structures



Regular expressions

- Describe regular grammars with (meta)characters
- Matched against texts to find instances of concepts
- Used in NLP particularly for (alpha)numeric entities

HW Henning Wachsmuth
Regular expressions
Ans: Henning Wachsmuth

Dear students,

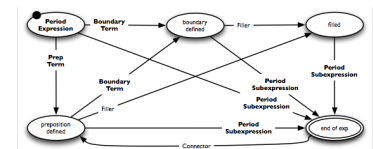
how does the mail tool infer that the following lines contain contact information:

Marty McFly
9303 Lyon Drive
Hill Valley, CA 95420
USA
marty.mcfly@deorian.bst

Does this have to do anything with regular expressions?

Benefits and limitations

- Effective for semi-closed classes, often very efficient
- Development complex for sophisticated concepts
- Restriction to sequential patterns limits applicability



References

Much content and many examples taken from

- Daniel Jurafsky and Christopher D. Manning (2016). Natural Language Processing. Lecture slides from the Stanford Coursera course.
<https://web.stanford.edu/~jurafsky/NLPCourseraSlides.html>.
- Daniel Jurafsky and James H. Martin (2009). Speech and Language Processing: An Introduction to Natural Language Processing, Speech Recognition, and Computational Linguistics. Prentice-Hall, 2nd edition.
- Friedhelm Meyer auf der Heide (2010). Einführung in Berechenbarkeit, Komplexität und Formale Sprachen. Begleitmaterial zur Vorlesung.
https://www.hni.uni-paderborn.de/fileadmin/Fachgruppen/Algorithmen/Lehre/Vorlesungsarchiv/WS_2009_10/Einfuehrung_in_die_Berechenbarkeit_K_u_f_S/skript.pdf
- Henning Wachsmuth (2015): Text Analysis Pipelines — Towards Ad-hoc Large-scale Text Mining. LNCS 9383, Springer.